

Computational Models - Exercise 5

19/1/08

1. Write a DFA for each of the following languages above the alphabet $\Sigma = \{a, b\}$:
 - a. Σ^*
 - b. \emptyset
 - c. $\{\varepsilon\}$
 - d. $\{w \mid |w| \bmod 3 = 0\}$
 - e. $\{w \mid w \text{ does not contain the substring 'abba'}\}$

2. For each of the following languages write a NFA and convert it to a DFA:
 - a. $\{w \mid w \text{ contains the substring 'aa' or doesn't contain the substring 'bab'}\}$
 - b. $\{w \in \{a, b\}^* \mid \text{There exists a partitioning of } w \text{ to } w = xy \text{ s.t. } x \text{ contains an even number of 'a's and } y \text{ contains an even number of 'b's}\}$

3. Write a regular expression for each of the following languages above $\Sigma = \{0, 1\}$:
 - a. $\{w \mid |w| \bmod 4 = 0\}$
 - b. $\{w \mid w \text{ contains exactly four '1's}\}$

4. For each of the following languages prove that if L is regular then L' is regular
 - a. $L' = \{xy \mid x \notin L \text{ and } y \in L\}$
 - b. $L' = \{x_2x_4 \dots x_{2n} \mid x_2, x_4, \dots, x_{2n} \in \Sigma, \exists x_1, x_3, x_{2n-1} \in \Sigma \text{ s.t. } x_1x_2 \dots x_{2n-1}x_{2n} \in L\}$, i.e. L' is the language of all strings that are a concatenation of the characters appearing in even indices in words of even length from L .

5. Prove if the following languages are regular:
 - a. $L_1 = \{w \mid w \in \{0, 1\}^*, w = w^R\}$
 - b. $\Sigma = \{0, 1, \#\}$, $L_2 = \{x\#y\#z \mid x, y, z \in \{0, 1\}^*, x + y = z\}$ (where x, y, z are numbers in binary representation).
 - c. $\Sigma = \{0, 1\}$, $L_3 = \{x_1y_1z_1 \dots x_ny_nz_n \mid x = x_n \dots x_1, y = y_n \dots y_1, z = z_n \dots z_1, x + y = z\}$, where x, y, z are numbers in binary representation and $x_1, \dots, x_n, y_1, \dots, y_n, z_1, \dots, z_n \in \Sigma$
 - d. $\Sigma = \{0, 1\}$, $L_4 = \{w \mid \text{the number of '01's in } w \text{ is equal to the number of '10's}\}$

6. Space complexity:
 - a. Prove that *Clique* is in *PSPACE* by giving an algorithm that solves *Clique* with polynomial space. ($\text{Clique} = \{\langle G, k \rangle \mid G \text{ is an undirected graph with a clique of size at least } k\}$)
 - b. Prove that *QSAT* (shown in class and recitation) is *NP-hard*. ($\text{QSAT} = \{\langle \phi \rangle \mid \phi \text{ is a true quantified Boolean formula}\}$)